

Markov property for Continuous time Processes

$$\left\{ \begin{aligned} &P(X(t) = j \mid X(s) = i, X(t_{n-1}) = i_{n-1}, \dots, X(t_1) = i_1) \\ &= P(X(t) = j \mid X(s) = i); \text{ where } 0 \leq t_1 \leq t_2 \leq \dots \leq t_{n-1} \leq s \leq t; \text{ and} \\ &i_1, i_2, \dots, i_{n-1}, i, j \in S \text{ are } (n+1) \text{ states in the} \\ &\text{state space; } \forall n \geq 1, n \in \mathbb{I}. \end{aligned} \right.$$

This is called Markov Property (in continuous time).

Defⁿ (Continuous time Markov chain or CTMC).

A continuous-time stochastic process $\{X(t) \mid t \geq 0\}$ is called a continuous-time Markov chain (CTMC) if it has the Markov property.

→ Memorylessness (Markov property).

→ time-homogeneity.

Defⁿ (Time Homogeneity).

We say that a CTMC is time-homogeneous if for any $s \leq t$ and any states $i, j \in S$

$$\begin{aligned} P(X(t) = j \mid X(s) = i) &= P(X(t-s) = j \mid X(0) = i) \\ &= P(X(t_1) = j \mid X(t+s+t_1) = i) \end{aligned}$$

the key thing to note & so on is that the difference in the time argument is $(t-s)$.

Not all CTMC need be time-homogeneous but in this course we will only consider time-homogeneous CTMC!

Meaning of time-homogeneity.

Whenever the process enters state i ; the way it evolves probabilistically from that pt is the same as if the process started in state i at time 0.

Defⁿ (Holding time) :- When the process enters state i , the time it spends there before it leaves state i is called the holding time.

T_i := holding time in state i
(just like in our motivating example of the n -server system).

Proposition :- $T_i \sim \text{exponential } D^n$.

Recall; $P_{i,j}$ = probability of going from state i to j

Similar to Discrete time Markov chain

$q_{i,j}$ = rate at which the system goes from state i to j
(rate at which the exponential alarm clocks go off).

Here, in CTMC, both $p_{i,j}$ & $q_{i,j}$ are fⁿ of time :- $p_{i,j}(t)$ & $q_{i,j}(t)$

$$p_{i,j} = \frac{q_{i,j}}{v_i} ; v_i = \sum_{j \in S} q_{i,j} < \infty$$

By defⁿ, $q_{i,i} = 0$

$$\Rightarrow q_{i,j} = v_i p_{i,j}$$

$(v_i = 0 \Rightarrow \text{state } i \text{ is absorbing state})$

The entries $p_{i,j}$ form a matrix $P = P(t)$ known as the Stochastic Matrix.

Note :- $q_{i,j}$ has more information about CTMC (continuous stochastic process) than $p_{i,j}$ b/c if we know all the $q_{i,j}$'s then we can find v_i & $p_{i,j}$. But if we know the $p_{i,j}$'s we cannot find $q_{i,j}$'s!

data we will use that the $q_{i,j}$'s from a matrix that is called the Generator matrix generates the $p_{i,j}$'s.

In many ways, $q_{i,j}$ are to CTMC what the $p_{i,j}$ are to DTMC.

$q_{i,j} > 0$ but $q_{i,j}$ need not be ≤ 1 like the $p_{i,j}$'s.

Stochastic Matrix

$P(t)$ comprise of $P_{ij}(t) = P(X(t)=j | X(0)=i)$

Note there is no "time step" in CTMC, instead we have $P_{ij}(t)$ which is a continuous ^{fn of} time. ~~tt~~

Often for CTMC; instead of writing $P_{ij}(t)$ we use upper case $P_{i,j}(t) \equiv P_{ij}(t)$

eg if the CTMC is a poisson process; then

$$P_{ij}(t) = P(\text{there are } j-i \text{ events (arrivals) in } t \text{ interval of time})$$
$$= \frac{(\lambda t)^{j-i} e^{-\lambda t}}{(j-i)!}$$

Chapman - Kolmogorov eqn for CTMC.

$$P_{ij}(t+s) = \sum_{k \in S} P_{kj}(s) P_{ik}(t)$$
$$\equiv \sum_{k \in S} P_{ik}(t) P_{kj}(s)$$

the above is the $(ij)^{\text{th}}$ entry of $P(t+s)$

$$|P(t+s) = P(t)P(s)|$$

Recall $T_i \sim \exp(\nu_i)$; $\nu_i = \sum q_{ij}$ Pg ③

$$f_{T_i}(h) = \nu_i e^{-\nu_i h}$$

$$\therefore P(T_i \leq h) = 1 - e^{-\nu_i h}$$

$$\Rightarrow P(T_i > h) = e^{-\nu_i h}$$

(comp. w/ n-server
motivating
example)

$$\begin{array}{c} \text{Taylor} \\ \text{expand} \\ \text{at } h=0 \end{array} 1 - \nu_i h + \frac{(\nu_i h)^2}{2!} - \dots$$

$$= 1 - \nu_i h + \underbrace{O(h)}_{\hookrightarrow O(h^2)}$$

$$\Rightarrow P(T_i \leq h) = 1 - (1 - \nu_i h + O(h)) \\ = \nu_i h + O(h)$$

$$\therefore P(0 \text{ transitions by time } h \mid X(0) = i) \\ = P(T_i > h) = 1 - \nu_i h + O(h)$$

$$P(\text{Exactly 1 transition by time } h \mid X(0) = i) \\ = 1 - P(0 \text{ transitions by time } h \mid X(0) = i) \\ = 1 - (1 - \nu_i h + O(h)) = \nu_i h + O(h)$$

& Likewise, it can be shown that - Try it out
yourself

$$P(2 \text{ or more transitions by time } h \mid X_0 = i) \\ = O(h)$$

Now

$$P(X(h) = j | X(0) = i) = (\underbrace{v_{ij}h}_{\text{exactly 1 event occurred}} + o(h)) p_{ij} + P(\underbrace{2 \text{ or more events}}_{p_{ij}})$$

$$= v_{ij} p_{ij} h + o(h) + o(h) \\ = v_{ij} p_{ij} h + o(h) \quad \text{--- (I)}$$

Similarly.

$$P(X(h) = i | X(0) = i) = 1 - v_{ii}h + o(h) + o(h) \\ = 1 - v_{ii}h + o(h) \quad \text{--- (II)}$$

Further

$$p_{ij}(t+h) = P(X(t+h) = j | X(0) = i)$$

law of

total

prob.

$$\sum_{k \in S} P(X(t+h) = j | X(h) = k, X(0) = i)$$

$$P(X(h) = k | X(0) = i)$$

Markov

prop.

$$\sum_{k \in S}$$

$$P(X(t+h) = j | X(h) = k) \cdot P(X(h) = k | X(0) = i)$$

reset-

clock

or

time homogeneity

$$\sum_{k \in S}$$

$$P(X(t+h) = j | X(0) = k) p_{ik}(h)$$

$$= \sum_{k \in S} p_{kj}(t) p_{ik}(h)$$

Basically Chapman Kolmogorov

Using (I) & (II) above.

$$p_{ij}(t+h) = p_{ij}(t) (1 - v_{ii}h + o(h)) + \sum_{k \neq i} p_{kj}(t) (v_{ik}h + o(h))$$

or equivalently

$$p_{ij}(t+h) - p_{ij}(t) = -v_i p_{ij}(t)h + \sum_{k \neq i} p_{kj}(t) v_i p_{ik}h + o(h)$$

Dividing by h ; taking limit $h \rightarrow 0$ & using $v_i p_{ik} = q_{ik}$

$$\boxed{p'_{ij}(t) = \sum_{k \neq i} q_{ik} p_{kj}(t) - v_i p_{ij}(t)} \quad \text{--- (A.1)}$$

or equivalently

$$\boxed{(P'(t))_{ij} = (Q P(t))_{ij}} \quad \text{--- (A.2)}$$

or

$$\boxed{P'(t) = Q P(t)} \quad \text{--- (A.3)}$$

(A.1), (A.2) & (A.3) are called Kolmogorov's Backward Eqns.

c) the infinitesimal generator matrix Q given by

$$\boxed{\begin{matrix} q_{ij} = q_{ij} & \forall i \neq j \\ q_{ii} = -v_i \end{matrix}} \quad \text{Identity matrix}$$

with bdy condⁿ $P(0) = I$

Likewise, we have Kolmogorov's Fwd eqns

$$P'(t) = P(t) Q \quad \text{--- (B.1)}$$

Kolmogorov's Bkwd & fwd eqns, w/ the
 bdy condn $P(0) = \mathbb{I}$, both have the
 same soln.

$$P(t) = e^{tQ} := \mathbb{I} + tQ + \frac{(tQ)^2}{2!} + \dots \quad (c)$$

∴ Even though we cannot normally
 obtain $P(t)$ in a simple closed form

We can use eqn (c) to obtain a
 numerical approxⁿ to $P(t)$ if $|s|$ is finite
 by truncating the ∞ sum to a finite
 sum.

* The solution $P(t) = e^{tQ}$ shows how
 basic the generator matrix Q is to
 the properties of CTMC.

We will find, in the subsequent
 lectures, that Q is also a key element
 for determining stng. Dⁿ of CTMC.

#

Stny Distributions.

pg(5)

Defⁿ:- Let $\{X(t) | t \geq 0\}$ be a CTMC w/ state space S , generator Q & matrixⁿ transition probabilityⁿ $TP(t)$

An $|S|$ -dimensional (row) vector $\vec{\pi} = (\pi_i)_{i \in S}$ w/ $\pi_i \geq 0 \forall i$ and

$\sum_{i \in S} \pi_i = 1$ is said to be a

Stationary Distribution if

$$\vec{\pi} = \vec{\pi} P(t) \quad \forall t \geq 0.$$

Question:- How does the generator Q relate to the definition of Stny.

Ans:- $\vec{\pi}$ is a D^n Stny Distribution $\Leftrightarrow \vec{\pi} = \vec{\pi} P(t), \forall t \geq 0$

$$P = e^{tQ} \quad \vec{\pi} = \sum_{n=0}^{\infty} \frac{(tQ)^n}{n!}, \forall t \geq 0$$

from prev. lecture on Kolmogorov eqns for CTMC

$$\Leftrightarrow \vec{\pi} - \vec{\pi} = \sum_{n=1}^{\infty} \frac{t^n}{n!} \vec{\pi} Q^n; \forall t \geq 0$$

$$\Leftrightarrow 0 = \vec{\pi} Q^n \quad \forall n \geq 1$$

$$\Leftrightarrow \boxed{\vec{\pi} Q = 0}$$

Sum of the terms is zero if and only if each summand is zero.

therefore the condition $\dot{\pi} = \pi P(t) \quad \forall t \geq 0$,
 that will be quite difficult to check,
 reduces to the much simpler condition

$$\boxed{\pi Q = 0}$$

↑
 a set of $|S|$
 linear eqns.

i.e. $\pi_j v_j = \sum_{i \neq j} \pi_i q_{ij}$

$$; v_j = \sum_{i \in S} q_{ji}$$

Physical interpretation.

long run
 proportion of time
 the process is in
 state j

rate of leaving
 state j when
 the process is
 in state j

long run rate of
 going from state
 i to state j

$$\Rightarrow \pi_j v_j = \text{long run rate of leaving state } j$$

$$\Rightarrow \sum_{i \neq j} \pi_i q_{ij} = \text{long run rate of going to state } j$$

$$\Rightarrow \text{"the long run rate out of state } j\text{"} = \text{"long run rate into state } j\text{"}$$

i.e. ~~that~~ it is a statement of dynamic equilibrium

& the eqn. $\pi Q = 0$ is also called
 the Global Balance eqn / Balance eqn

GLOBAL BALANCE EQNS

Detailed Balance eqns. (Also known as Local Balance) Pg (6)

For a continuous time Markov Ch (CTMC) w/ transition matrix Q ; if π_i can be found s.t. for every pair of states i & j

$$\pi_i Q_{ij} = \pi_j Q_{ji} \quad \text{--- (Detailed Balance condition)}$$

holds; then by summing over j , the global balance eqns are satisfied & $\vec{\pi}$ is stng. D^n of the process.

* If such a solution can be found the resulting eqns are usually much easier than directly solving the global balance eqns.

☞ A CTMC is reversible (\Leftrightarrow) detailed balance holds for every (i, j)

Note :- The equivalent detailed balance for DTMC is

$$\pi_i P_{ij} = \pi_j P_{ji} \quad \forall i, j \text{ pairs.}$$

Limiting probabilities -

for a CTMC $\{X(t) | t \geq 0\}$,
 $\lim_{t \rightarrow \infty} P_{ij}(t) = \lim_{t \rightarrow \infty} P(X(t) = j | X(0) = i)$

$$\equiv \pi_j$$

i.e. limiting probability \equiv Stnary prob.
 $(t \rightarrow \infty)$ D^n .

Application of Local (detailed) Balance Equations

Local balance :- $\pi_i q_{ij} = \pi_j q_{ji} \quad \forall i, j \in D^n$
 $i \neq j$

there are $|D|^2$ such eqns.

but typically most of the eqns are trivially satisfied b/c
 $q_{ij} = q_{ji} = 0$

* Recall that for global balance (& Stnary Cond) to exist; it is not necessary that local balance condn. ^{always} holds but if local balance does hold then surely global balance (& Stnary) ~~do~~ hold.

* One quick way to check if local balance does **NOT** hold for stationarity is to check if there are any rates

q_{ij} and q_{ji} s.t. $q_{ij} > 0$ and pg(7)

$q_{ji} = 0$ or $q_{ij} = 0$ & $q_{ji} > 0$.

(In these cases, the local balance
route to investigating stationarity
will be futile!)
